

# Programming Abstractions

**Week 8-2: MiniScheme C**

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# What can MiniScheme do at this point?

MiniScheme B has constant numbers

MiniScheme B has pre-bound symbols that are in the `init-env`

# Recall

`(parse input)` — Parses the input, at this point either a number or a variable, and returns a `(lit-exp num)` or `(var-exp sym)`

`(eval-exp tree e)` — Evaluates the parse tree in the environment `e`, returning a value

What does `(parse 15)` return?

A. 15

B. `'(lit-exp 15)`

C. It's an error

What does `(parse 'z)` return?

A. `'(lit-exp z)`

B. `'(var-exp z)`

C. It's an error

What does `(eval-exp (var-exp 'z) environment)` do?

- A. Returns what `z` is bound to in `environment`
- B. It's an error
- C. It looks up with `z` is bound to, returning the result or causing an error if `z` is not bound
- D. Something else

What does `(eval-exp (lit-exp 108) environment)` do?

- A. Returns what 108 is bound to in environment
- B. It's an error
- C. It looks up with 108 is bound to, returning the result or causing an error if 108 is not bound
- D. Returns 108
- E. Something else

# Homeworks 6 and 7

Multiple steps, each adding parts to the MiniScheme interpreter

For each new type of expression

- ▶ Add a new data type
  - `if-exp`
  - `let-exp`
  - etc.
- ▶ Add constructors, recognizers and accessors
- ▶ Modify `parse` to produce those
- ▶ Modify `eval-exp` to interpret them

```
EXP → number
      | symbol
      | ( if EXP EXP EXP )
      | ( let ( LET-BINDINGS ) EXP )
      | ( letrec ( LET-BINDINGS ) EXP )
      | ( lambda ( PARAMS ) EXP )
      | ( set! symbol EXP )
      | ( begin EXP* )
      | ( EXP EXP* )
LET-BINDINGS → LET-BINDING*
LET-BINDING → [ symbol EXP ]
PARAMS → symbol*
```



# Let's add arithmetic and some list procedures

## MiniScheme C

Let's add +, −, \*, /, car, cdr, cons, etc.

Students find this to be the hardest part of the project

- ▶ It's the first complex part
- ▶ It contains some things that make more sense later, once we add lambda expressions

# Many ways to call procedures

```
(+ 2 3)
```

```
((lambda (x y) (+ x y)) 2 3)
```

```
(let ([f +]) (f 2 3))
```

The parser can't identify primitive procedures like + because symbols like f may be bound to primitive procedures

- ▶ It can't tell because the parser **does not have access to the environment**

All that the parser can do is recognize a procedure application and parse

- ▶ the procedure; and
- ▶ the arguments

# Enter lists

So far, the input to MiniScheme A and B has just been a number or a symbol

If the input is a list, then the kind of expression it represents depends on the first element

- ▶ If the first element is 'lambda, it's a lambda expression
- ▶ If the first element is 'let, it's a let expression
- ▶ If the first element is 'if, it's an if-then-else expression
- ▶ etc.

Applications don't have keywords, so any nonempty list for which the first element is not one of our supported keywords is an application

# Procedure applications

## MiniScheme C

$EXP \rightarrow$  number                    parse into `lit-exp`  
          | symbol                    parse into `var-exp`  
          | (*EXP EXP\**)            parse into `app-exp`

An `app-exp` is a new data type that stores

- ▶ The parse tree for a procedure
- ▶ A list of parse trees for the arguments

Procedures to implement

- ▶ `(app-exp proc args)`
- ▶ `(app-exp? exp)`
- ▶ `(app-exp-proc exp)`
- ▶ `(app-exp-args exp)`

# Recursive implementation

## Parsing

Expressions are recursive:  $EXP \rightarrow ( EXP EXP^* )$

When parsing an application expression, you want to parse the sub expressions using parse

```
(define (parse input)
  (cond [(number? input) (lit-exp input)]
        [(symbol? input) (var-exp input)]
        [(list? input)
         (cond [(empty? input) (error ...)]
               [else (app-exp (parse (first input))
                               (...))])]
        [else (error 'parse "Invalid syntax ~s" input)]))
```

Parse the procedure

Parse the arguments

# How should you parse the arguments?

Consider input that looks like  
`((lambda (x y) x) 2 3)` or  
`(f 4 5 6)`

The procedure part can be parsed with `(parse (first input))`

How should you parse the arguments?

# Evaluating an `app-exp`

Evaluate the procedure part

Evaluate each of the arguments

If the procedure part evaluates to a primitive procedure, call a procedure you'll write that will perform the operation on the arguments

▶ E.g., if the primitive procedure is `*`, then you'll want to call `*` on the arguments

The tricky part is what does it mean to evaluate the procedure part?

# Evaluating the procedure part of an `app-exp`

Consider the input `'(+ 2 3 4)`

The procedure part is `' +` which will be parsed as `'(var-exp +)`

Variable reference expressions are evaluated by looking the symbol up in the current environment

Therefore, we need our initial environment to contain a binding for the symbol `' +` (and all the other primitive procedures we want to support)



# prim-proc data type

We can create a new data type prim-proc

- `(prim-proc symbol)`
- `(prim-proc? value)`
- `(prim-proc-symbol value)`

The prim-proc is *only* used to interpret expressions so where should this data type and its procedures be defined?

# Adding primitives to our initial environment

```
(define primitive-operators  
  '(+ - * /))
```

```
(define prim-env  
  (env primitive-operators  
        (map prim-proc primitive-operators)  
        empty-env))
```


```
(define init-env  
  (env '(x y) '(23 42) prim-env))
```

# Evaluating an `app-exp`

Recall: `app-exp` stores the parse tree for the procedure and a list of parse trees for the arguments

We need to evaluate all of those; add something like the following to `eval-exp`

```
[ (app-exp? tree)
  (let ([proc (eval-exp (app-exp-proc tree) e)]
        [args ...])
    (apply-proc proc args)) ]
```



# Applying a procedure

The `apply-proc` procedure takes an evaluated procedure and a list of evaluated arguments

It can look at the procedure and determine if it's a primitive procedure

- ▶ If so, it will call `apply-primitive-op`
- ▶ If not, it's an error for now; later, we'll add code to deal with non-primitive procedure (i.e., normal lambdas)

```
(define (apply-proc proc args)
  (cond [(prim-proc? proc)
        (apply-primitive-op (prim-proc-symbol proc) args)]
        [else (error 'apply-proc "Bad proc: ~s" proc)]))
```

# Applying primitive operations

## (apply-primitive-op op args)

apply-primitive-op takes a symbol (such as '+' or '\*') and a list of arguments

You probably want something like

```
(define (apply-primitive-op op args)
  (cond [(eq? op '+) (apply + args)]
        [(eq? op '*) (apply * args)]
        ...
        [else (error "...)]))
```

What is returned by `(parse '(* 2 3))`?

A. `'( (prim-proc *) 2 3 )`

B. `'( (prim-proc *) (lit-exp 2) (lit-exp 3) )`

C. `'( app-exp (prim-proc *) ((lit-exp 2) (lit-exp 3)) )`

D. `'( var-exp * (lit-exp 2) (lit-exp 3) )`

E. `'( app-exp (var-exp *) ((lit-exp 2) (lit-exp 3)) )`

When evaluating an `app-exp`, the procedure and each of the arguments are evaluated. For example, when evaluating the result of `(parse '(- 20 5))`, there will be three recursive calls to `eval-exp`, the first of which is evaluating `(var-exp '-)`.

What is the result of evaluating `(var-exp '-)`?

- A. `#<procedure:->` (i.e., the procedure - itself)
- B. `'(app-exp -)`
- C. `'(prim-proc -)`
- D. It's an error because `-` requires arguments

What is the result of `(eval-exp (parse '(* 4 5)) init-env)`?

A. 20

B. `'(app-exp (var-exp *) ((lit-exp 4) (lit-exp 5)))`

C. `'(prim-proc * 4 5)`

D. `'(prim-proc (var-exp *) (lit-exp 4) (lit-exp 5))`

E. `'(app-exp (prim-proc *) 4 5)`



# Why go to all that trouble?

In a later version of MiniScheme, we'll implement lambda

We'll deal with this by adding a line to `apply-proc` that will apply closures

# Adding other primitive procedures

In addition (pardon the pun) to `+`, `-`, `*`, and `/`, you'll add several other primitive procedures

- ▶ `add1`
- ▶ `sub1`
- ▶ `negate`
- ▶ `list`
- ▶ `cons`
- ▶ `car`
- ▶ `cdr`

And you'll add a new variable `null` bound to the empty list

**What does `(car (list 3 5 2))` parse to?**

**What does `(car (list 3 5 2))` parse to?**

```
'(app-exp (var-exp car)
          ((app-exp (var-exp list)
                    ((lit-exp 3)
                     (lit-exp 5)
                     (lit-exp 2))))))
```

# Adding additional primitive procedures

1. Add the procedure name to `primitive-operators`
2. Add a corresponding line to the `cond` in `apply-primitive-op`

E.g.,

```
[(eq? op 'car) (car (first args))]
```

```
[(eq? op 'list) args]
```

# What can MiniScheme C do?

Numbers

Pre-defined variables

Procedure calls to built-in procedures